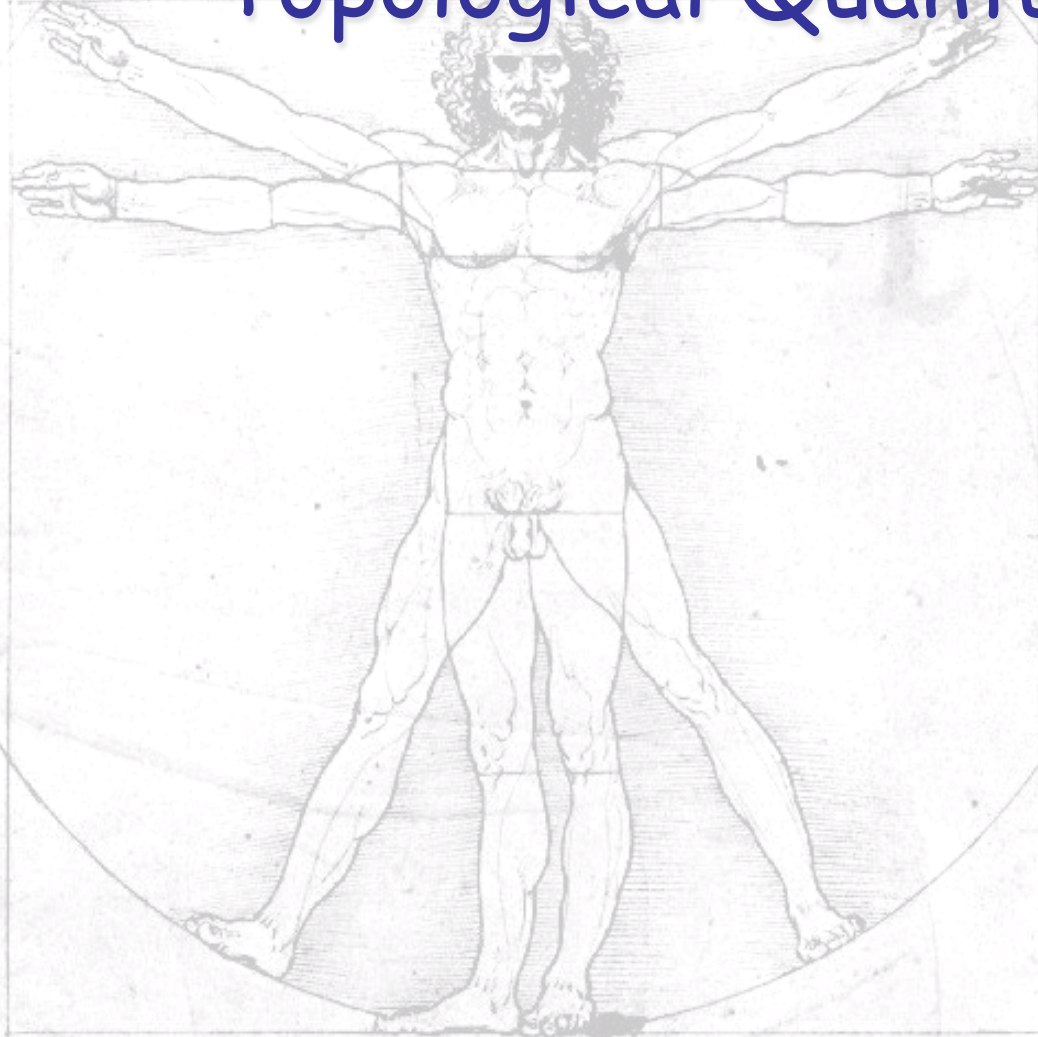


# Topological Quantum Computation

Jiannis K. Pachos

Introduction



*Bertinoro, June 2013*

**EPSRC**

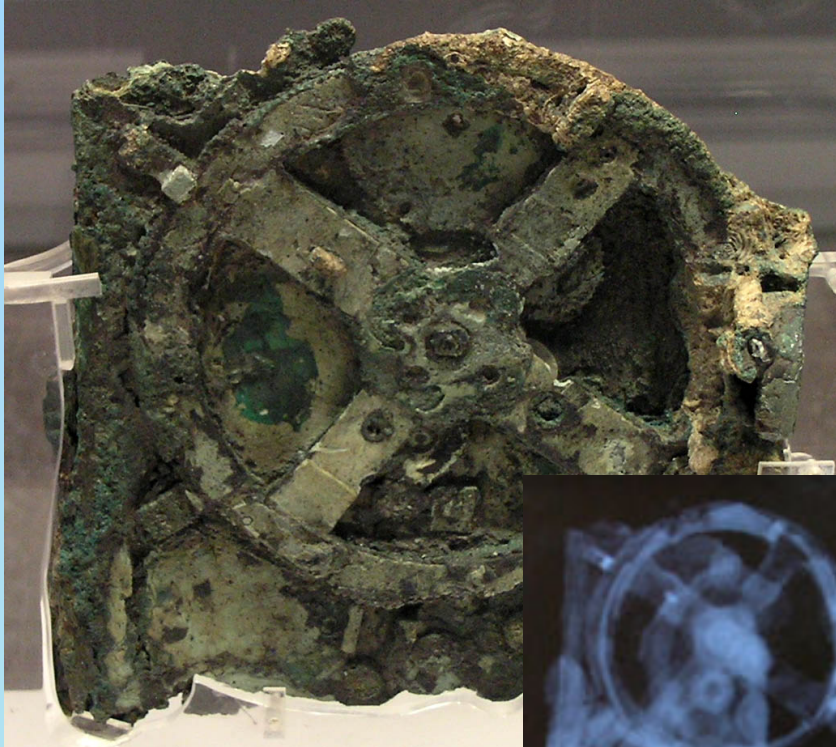
Engineering and Physical Sciences  
Research Council



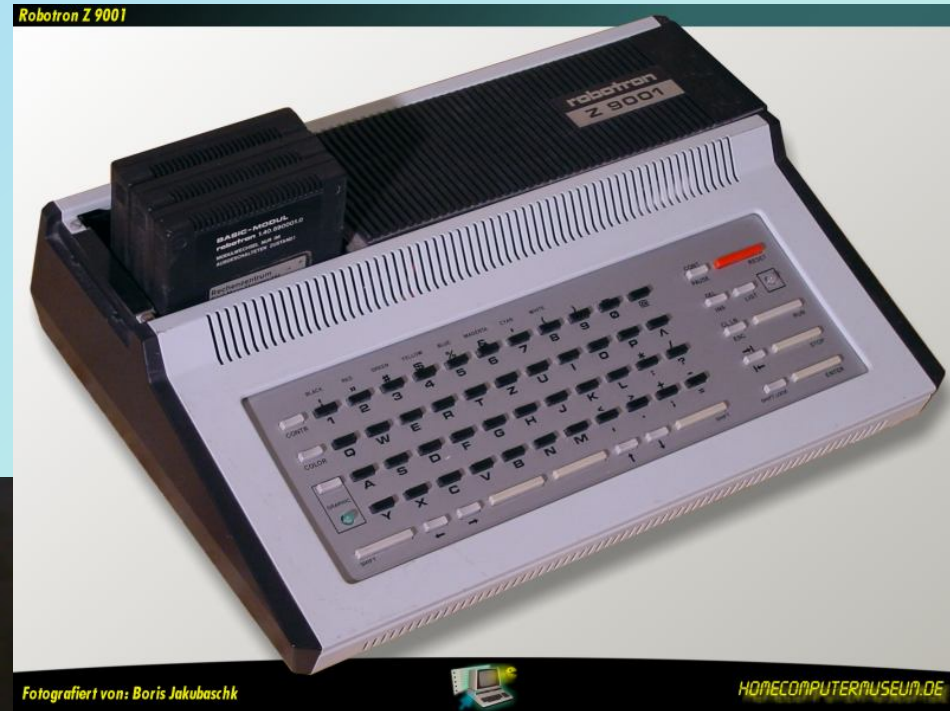
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# Computers

Antikythera mechanism



Robotron Z 9001



Analogue computer

Digital computer: 0 & 1

# Quantum computers

- Quantum binary states  $|0\rangle, |1\rangle$  (spin,...)

$$|\Psi\rangle = a|0\rangle + b|1\rangle \quad (\text{quantum superposition})$$

$$|a|^2 + |b|^2 = 1$$

If you **measure**  $|\Psi\rangle$  then  $P_0 = |a|^2, P_1 = |b|^2$

This is the quantum bit or “qubit”

- Many qubits (two):

$$|\Psi\rangle = |0\rangle \otimes |0\rangle \quad \text{product state (classical)}$$

$$|\Psi\rangle = a|0\rangle \otimes |0\rangle + b|1\rangle \otimes |1\rangle \quad \text{entangled state}$$

# Quantum computers

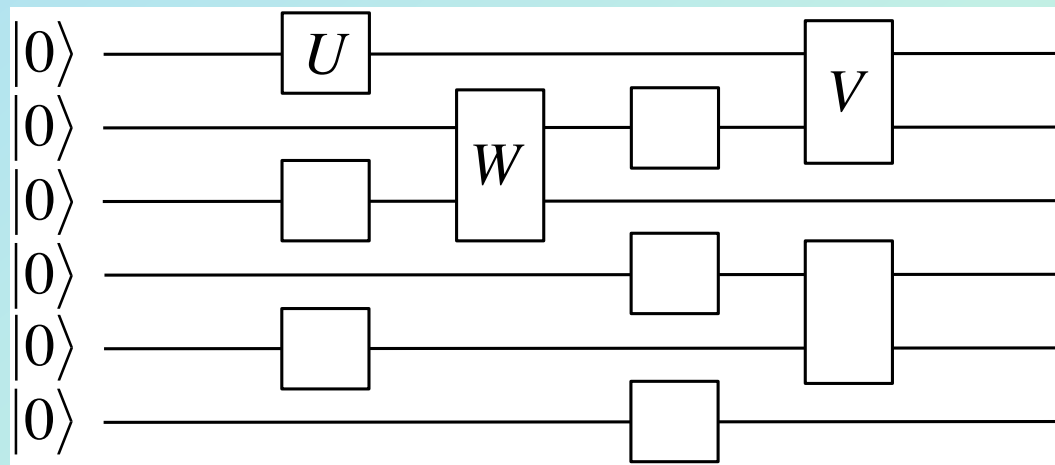
- Quantum binary gates  $\rightarrow$  **Unitary matrices**

$$U_{1q}|0\rangle = a|0\rangle + b|1\rangle$$

$$U_{2q}|0\rangle \otimes |0\rangle = a|0\rangle \otimes |0\rangle + b|1\rangle \otimes |1\rangle$$

- Universality:** Any  $U_{1q}$  and a  $U_{2q} \Rightarrow$  any q. algorithm.
- Quantum circuit model:

Input  
(product state)  
Initiate



Measure  
(product state)  
Output

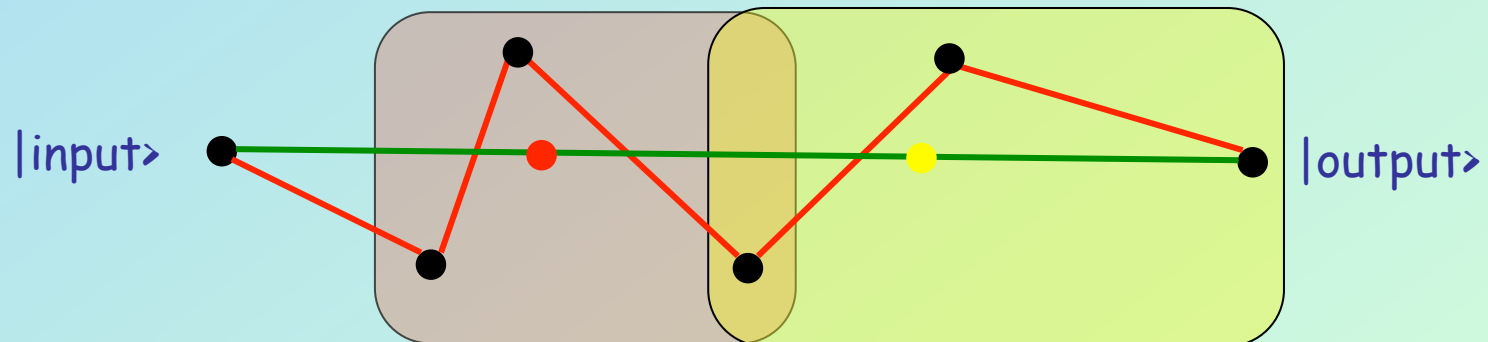


# Quantum computers

- Quantum Algorithm:



$$|\Psi_{\text{output}}\rangle = U_n \dots U_2 U_1 |\Psi_{\text{input}}\rangle$$



# Quantum computers: Why?

- **Computational complexity**

Problems that can be solved in:

- polynomial time (easy)

- exponential time (hard)

as a function of input size.

- **Classical computers:**

**P:** polynomially easy to solve

**NP:** polynomially easy to verify solution

- **BQP:** polynomially easy to solve with QC

# Quantum computers: Why?

- Factoring (Shor)

```
18070820886874048059516561644059055662781025167694013491701270214
50056662540244048387341127590812303371781887966563182013214880557
=
39685999459597454290161126162883786067576449112810064832555157243
×
45534498646735972188403686897274408864356301263205069600999044599
```

quantum hackers *exponentially* better than classical hackers!

- Searching objects (Grover): where is ♣?

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# Quantum computers: Why?

- **Factoring algorithm (Shor):**
  - Exponentially faster than known classical algorithm, but we do not know if there is a better classical one...
- **Searching algorithm (Grover):**
  - Quadratic speed up (optimal),  
does not change complexity class...
- Still important enough: worth investigating...
- **Errors during QC are too catastrophic.**



# Topological quantum computers: Why?

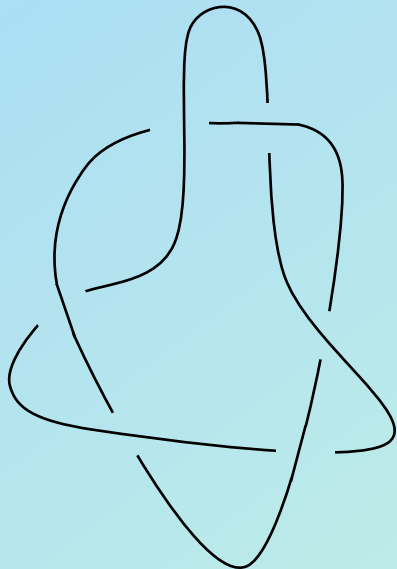


Topology promises to solve the problem of **errors** that inhibit the experimental realisation of quantum computers...

...and it is a lot of fun :-)

# Geometry - Topology

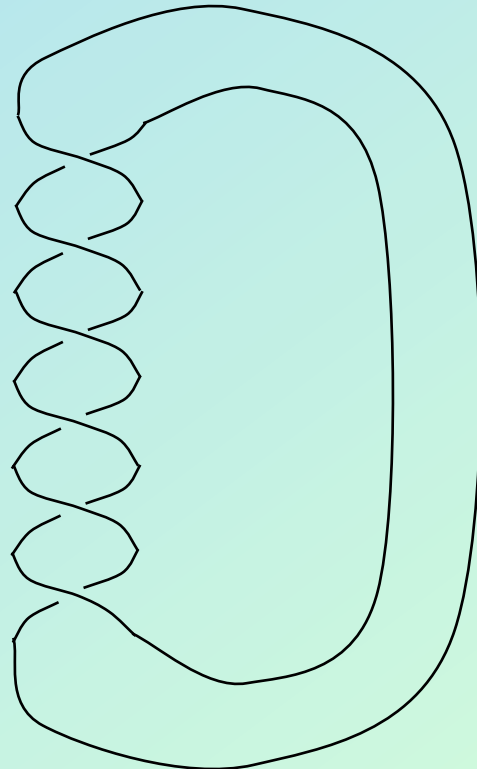
- **Geometry**
  - Local properties of object
- **Topology**
  - Global properties of object



geom.

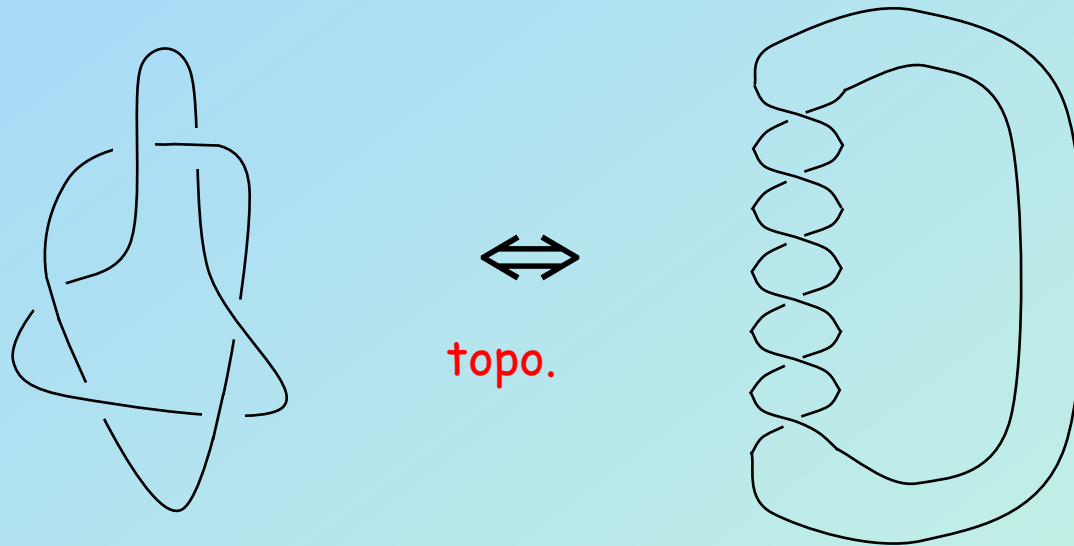


topo.



# Topology of knots and links

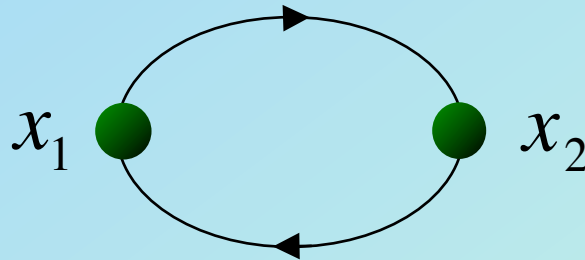
Are two knots equivalent?



- Algorithms exist from the '60s
- **Extremely** time consuming...
- Common problem (speech recognition, ...)
- *Mathematically **Jones polynomials** can recognise if two knots are inequivalent.*

# Particle statistics

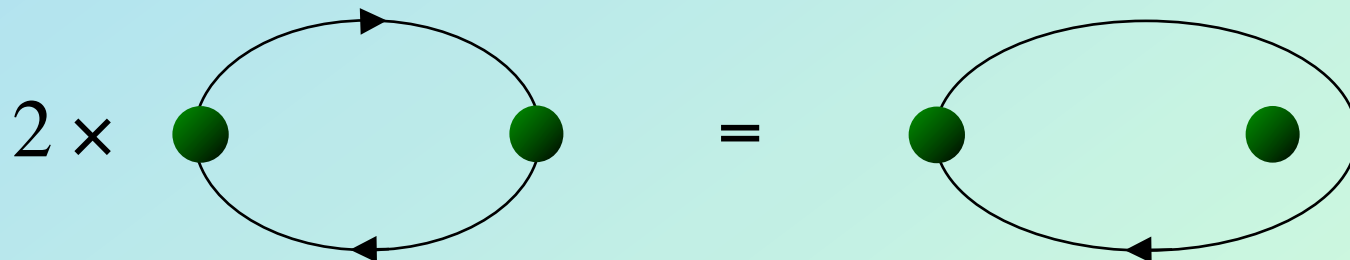
Exchange two identical particles:



Statistical symmetry:

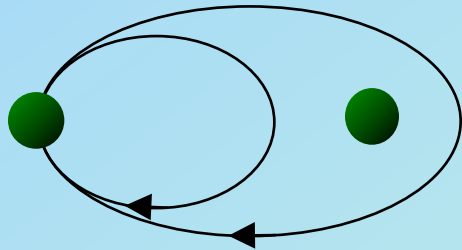
Physics stays the same, but  $|\Psi\rangle$  could change!

$$|\Psi(x_1, x_2)\rangle = ??? |\Psi(x_2, x_1)\rangle$$



# Anyons and statistics

3D



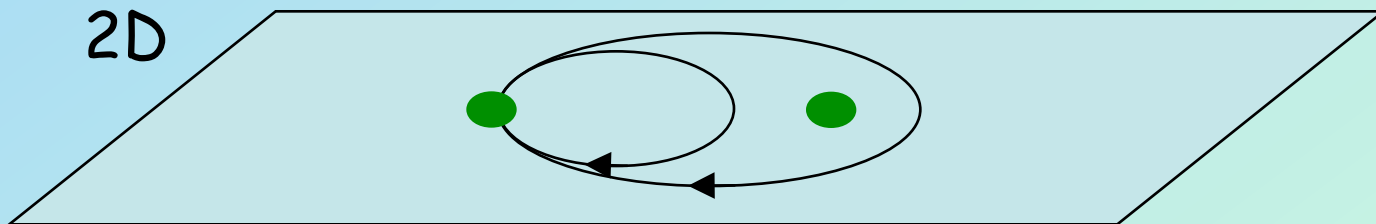
Bosons

$$|\Psi\rangle \rightarrow |\Psi\rangle$$

Fermions

$$|\Psi\rangle \rightarrow e^{i2\pi} |\Psi\rangle$$

2D



$$|\Psi\rangle \rightarrow e^{i2\phi} |\Psi\rangle$$

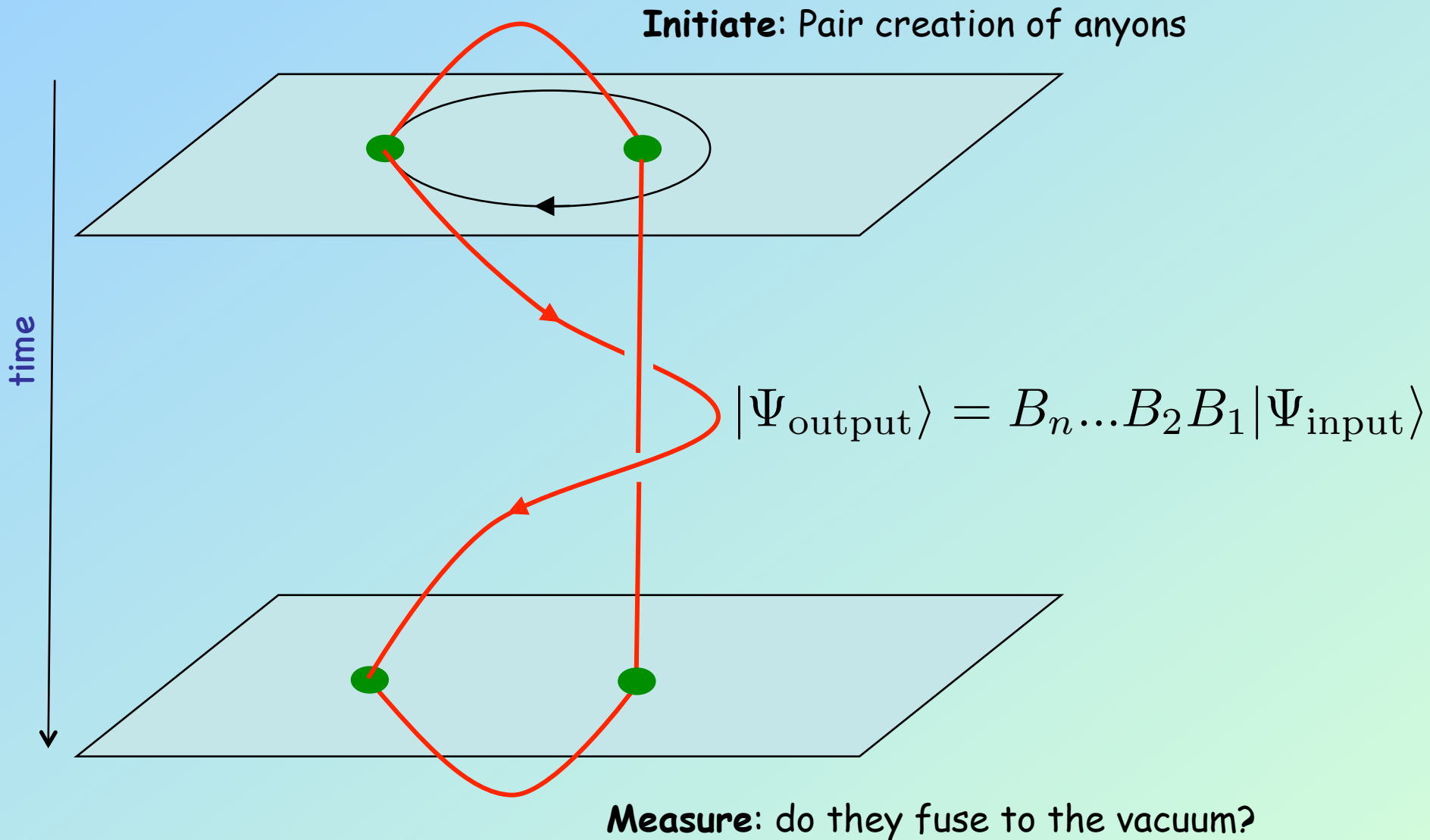
$$|\Psi\rangle \rightarrow B |\Psi\rangle$$

**Anyons**

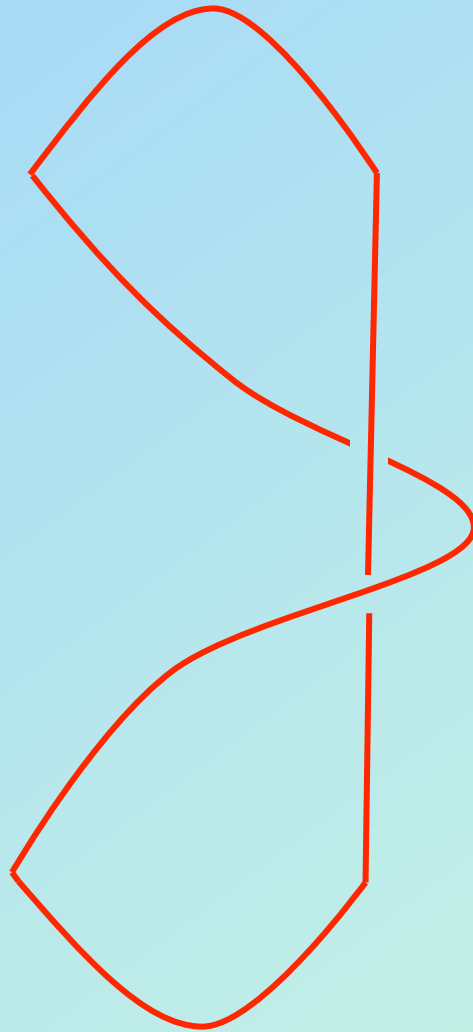
Like a  
quantum gate



# Anyons, statistics and knots



# Anyons, statistics and knots

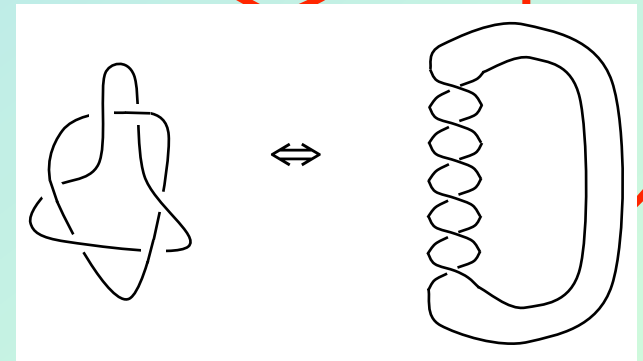
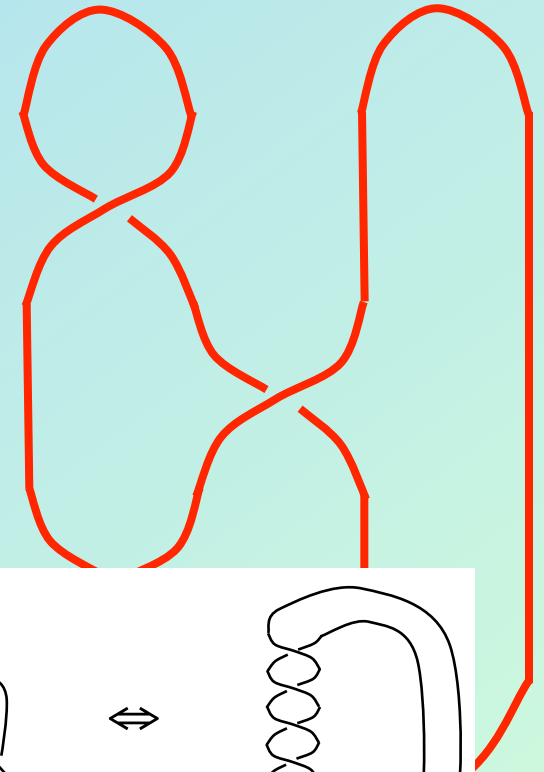


$$|\Psi_{\text{output}}\rangle = B_n \dots B_2 B_1 |\Psi_{\text{input}}\rangle$$

# Anyons and knots

Assume I can generate anyons in the laboratory.

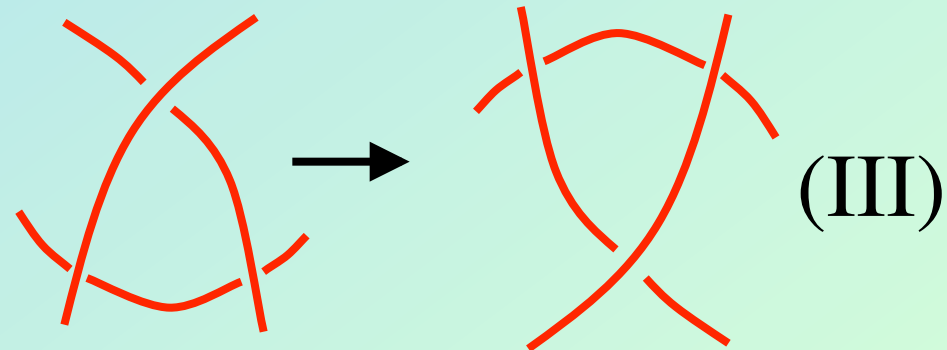
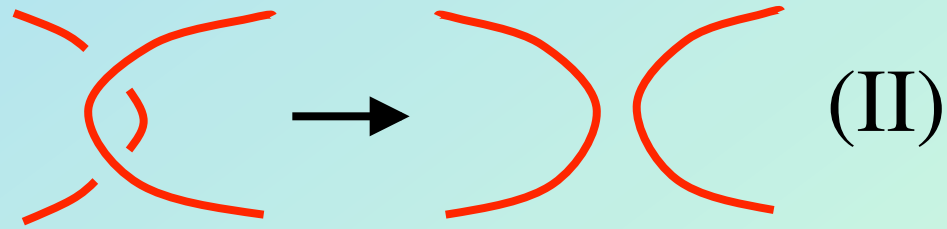
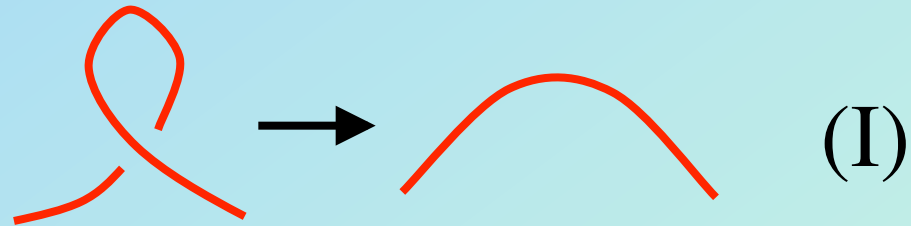
- The state of anyons is efficiently described by their **world lines**.
- Creation, braiding, fusion.
- The final **quantum state** of anyons is **invariant** under **continuous** deformations of strands.



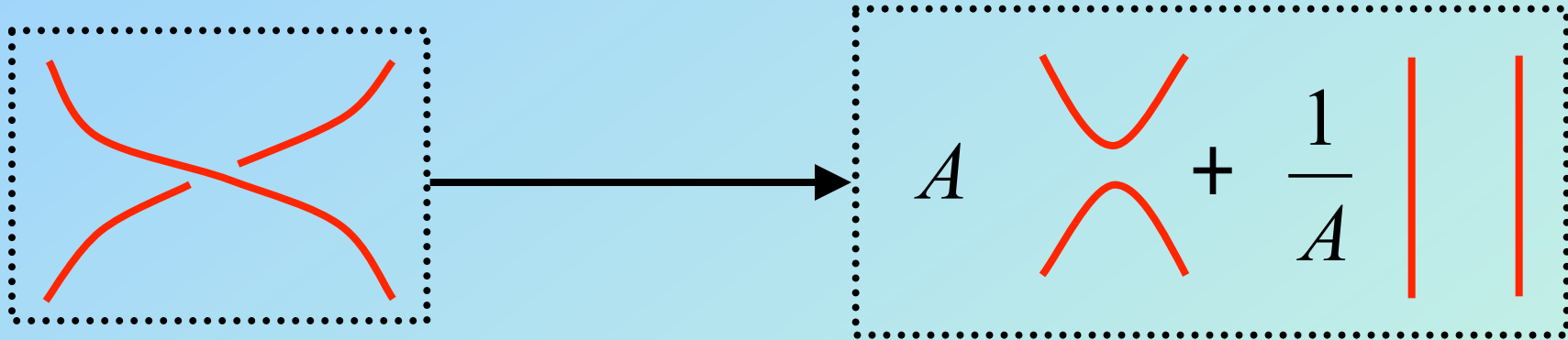
# The Reidemeister moves

**Theorem:**

Two knots can be deformed continuously one into the other iff one knot can be transformed into the other by local moves:

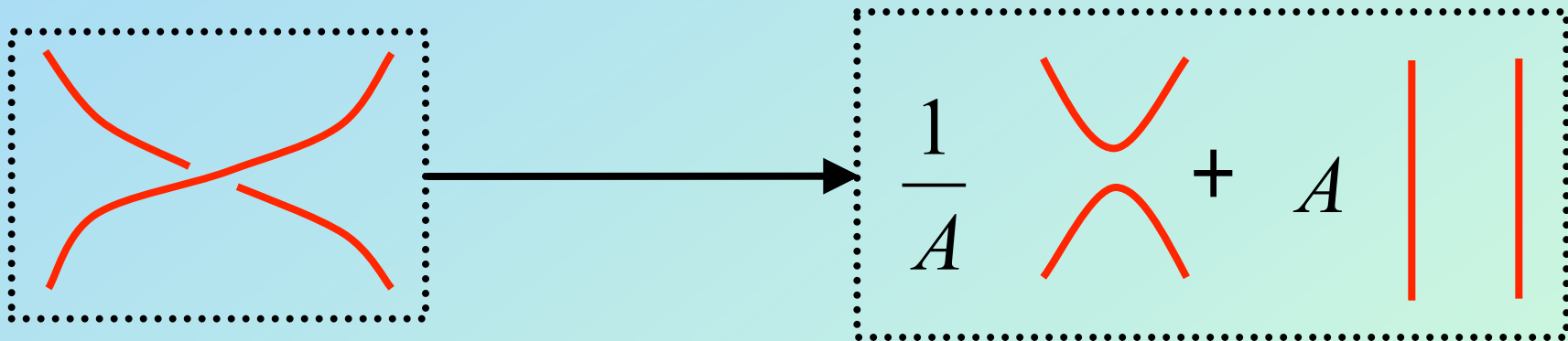


# Skein relations



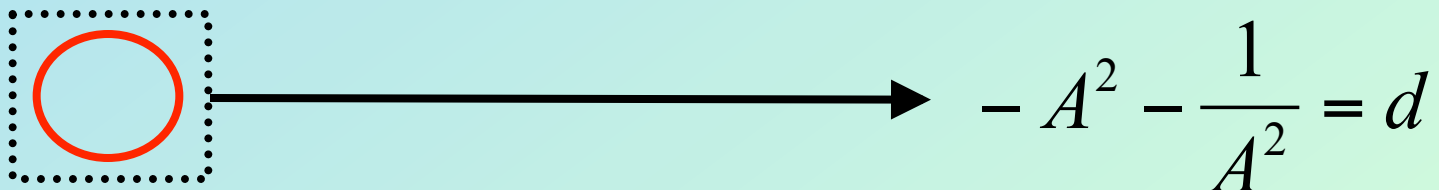
A diagram showing a crossing of two red strands in a dotted box. An arrow points to the right, where the result is shown in another dotted box. The result is the sum of two terms: the first term is the coefficient  $A$  multiplied by a diagram of two red strands meeting at a cusp (a V-shape pointing up and a V-shape pointing down), and the second term is the coefficient  $\frac{1}{A}$  multiplied by a diagram of two parallel vertical red strands.

$$A \begin{array}{c} \text{V-shape} \\ \text{V-shape} \end{array} + \frac{1}{A} \begin{array}{|c|} \hline \\ \hline \end{array} \begin{array}{|c|} \hline \\ \hline \end{array}$$



A diagram showing a crossing of two red strands in a dotted box. An arrow points to the right, where the result is shown in another dotted box. The result is the sum of two terms: the first term is the coefficient  $\frac{1}{A}$  multiplied by a diagram of two red strands meeting at a cusp (a V-shape pointing up and a V-shape pointing down), and the second term is the coefficient  $A$  multiplied by a diagram of two parallel vertical red strands.

$$\frac{1}{A} \begin{array}{c} \text{V-shape} \\ \text{V-shape} \end{array} + A \begin{array}{|c|} \hline \\ \hline \end{array} \begin{array}{|c|} \hline \\ \hline \end{array}$$



A diagram showing a single red circle in a dotted box. An arrow points to the right, where the result is given as an equation:  $-A^2 - \frac{1}{A^2} = d$ .

$$-A^2 - \frac{1}{A^2} = d$$



# Skein and Reidemeister

$$\begin{aligned}
 & \text{Diagram 1} = \frac{1}{A} \text{Diagram 2} + A \text{Diagram 3} = \\
 & = \text{Diagram 4} + \frac{1}{A^2} \text{Diagram 5} + A^2 \text{Diagram 6} + \text{Diagram 7}
 \end{aligned}$$

Diagram 1: A red line with a loop and a crossing.

Diagram 2: A red line with a loop and a crossing.

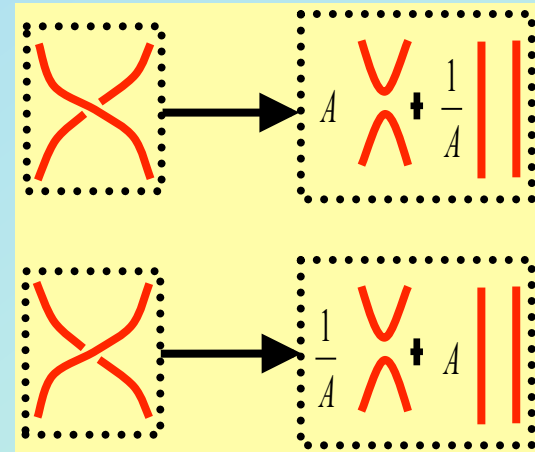
Diagram 3: A red line with a crossing.

Diagram 4: A red line with a loop and a crossing, with a black line passing through the loop.

Diagram 5: A red line with a crossing, with a black line passing through the crossing.

Diagram 6: A red line with a crossing, with a black line passing through the crossing.

Diagram 7: A red line with a crossing.

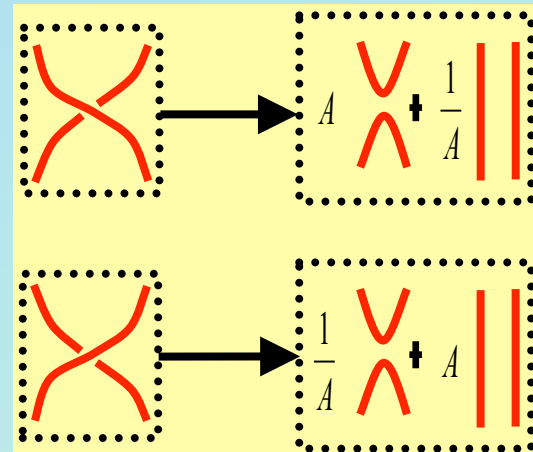


Reidemeister move (II) is satisfied. Similarly (III).

# Kauffman bracket

The Skein relations give rise to the **Kauffman bracket**:

$$\text{Skein}(\text{red knot}) = \langle L \rangle(A)$$



$$\langle \text{figure-eight} \rangle = A \langle \text{circle} \rangle + A^{-1} \langle \text{two circles} \rangle = A + dA^{-1} = (-A)^{-3}$$

$$\langle \text{trefoil} \rangle = A \langle \text{two concentric circles} \rangle + A^{-1} \langle \text{circle} \rangle = Ad + A^{-1} = (-A)^3$$

$$\langle \text{link} \rangle = A \langle \text{trefoil} \rangle + A^{-1} \langle \text{figure-eight} \rangle = -A^4 - A^{-4}$$

# Jones polynomial

The Skein relations give rise to the **Kauffman bracket**:

$$\text{Skein}(\text{link}) = \langle L \rangle(A)$$

To satisfy **move (I)** one needs to define **Jones polynomial**:

$$V_L(A) = (-A)^{3w(L)} \langle L \rangle(A)$$

$w(L)$  is the writhe of link. Easily computable.

# Jones polynomials

- If two links have different Jones polynomials then they are inequivalent

=> use it to distinguish links

- Jones polynomials keep:

only topological information, no geometrical

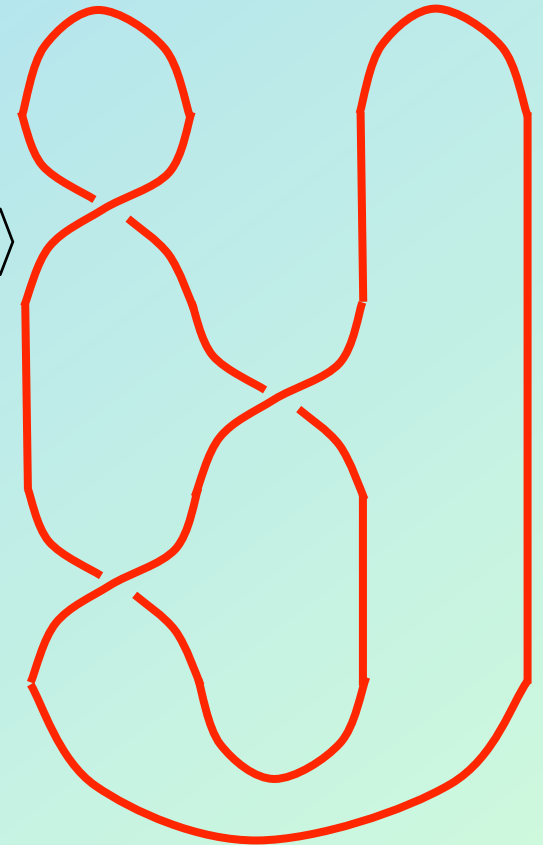
# Jones polynomial from anyons

Braiding evolutions of anyonic states:

$$|\Psi_{\text{final}}\rangle = B_n \dots B_2 B_1 |\Psi_{\text{initial}}\rangle$$

$$\begin{aligned} \langle \Psi_{\text{initial}} | \Psi_{\text{final}} \rangle &= \langle \Psi_{\text{initial}} | B_n \dots B_2 B_1 | \Psi_{\text{initial}} \rangle \\ &= \frac{1}{d^{n/2-1}} \langle L(B) \rangle \end{aligned}$$

- *Simulate the knot with braiding anyons*
- Translate it to circuit model:  
     $\Leftrightarrow$  find trace of matrices





# Jones polynomial from QC

Evaluating Jones polynomials is a #P-hard problem.

Belongs to BQP class.

With quantum computers it is **polynomially** easy to *approximate* with additive error.

[Freedman, Kitaev, Larsen, Wang (2002);  
Aharonov, Jones, Landau (2005);  
*et al.* Glaser (2009);  
Kuperberg (2009)]

# Conclusions

Jones polynomials are used for quantum applications:

- **encrypt** quantum information
- **quantum money**
- ...

Topological systems that can support **anyons** are currently **engineered...**

<http://quantum.leeds.ac.uk/~jiannis>

